



June 29, 2017



Exascale Scalability

• Which scalability metrics are important for Exascale?

- Performance (obviously!)
- What are the contributing factors?

• How can we demonstrate these principles today?

• Our architectural vision needs qualification



Second Thing First... Qualification



- How can we demonstrate and qualify scalability metrics for supercomputing?
- NNSA ASC Advanced Technology Platform
 - LANL / SNL Trinity Supercomputer

We already build BIG...



LANL / SNL Trinity System

Rank	Site	System	Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	National Supercomputing Center in Wuxi China	Sunway TaihuLight - Sunway MPP, Sunway SW26010 260C 1.45GHz, Sunway NRCPC	10,649,600	93,014.6	125,435.9	15,371
2	National Super Computer Center in Guangzhou China	Tianhe-2 (MilkyWay-2) - TH-IVB-FEP Cluster, Intel Xeon E5-2692 12C 2.200GHz, TH Express- 2, Intel Xeon Phi 31S1P NUDT	3,120,000	33,862.7	54,902.4	17,808
3	Swiss National Supercomputing Centre (CSCS) Switzerland	Piz Daint - Cray XC50, Xeon E5-2690v3 12C 2.6GHz, Aries interconnect , NVIDIA Tesla P100 Cray Inc.	361,760	19,590.0	25,326.3	2,272
10	DOE/NNSA/LANL/SNL United States	Trinity - Cray XC40, Xeon E5-2698v3 16C 2.3GHz, Aries interconnect Cray Inc.	301,056	8,100.9	11,078.9	4,233

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Why Trinity?

- First instantiation of NNSA ASC Advanced Technology Platform
 - Establishes foundation for Exascale
 - Meet future needs of current applications
 - Enable adaptation to new methodologies





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Trinity Architecture Overview



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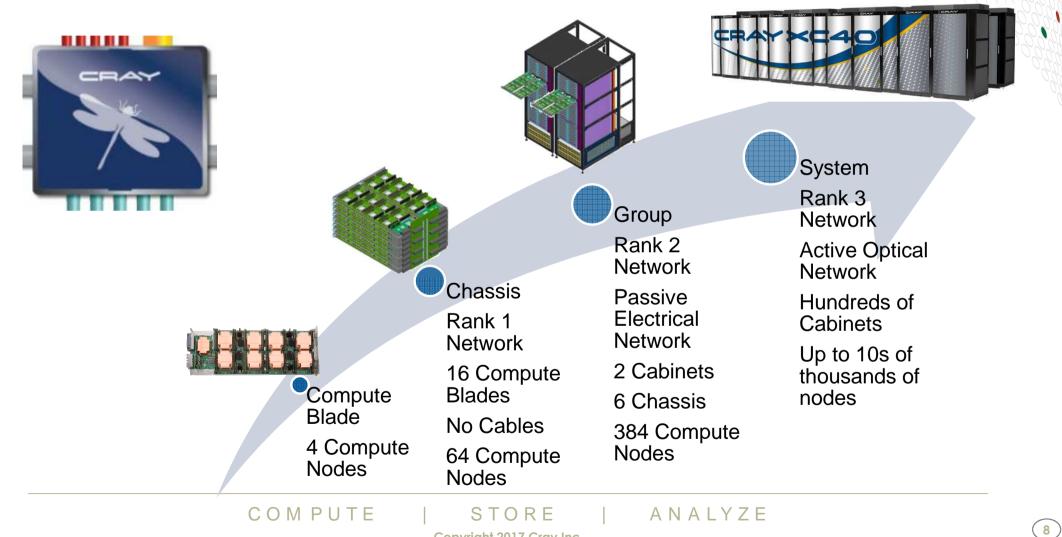
Trinity							
Node architecture	KNL & Haswell	Haswell Partition	KNL Partition				
Memory capacity	2.11 PB	>1 PB	>1 PB				
Memory BW	>7 PB/s	>1 PB/s	>1 PB/s				
Peak Flops	42.2 PE	11.5 PF	30.7 PF				
Number of nodes	19,000+	>9,500	9,500				
Number of cores	>760,000	>190,000	>570,000				
PFS capacity	>80 PB						
Burst Buffer capacity	3.7 PB						
Network interconnect	Cray Aries Network						
Number of cabinets	112						



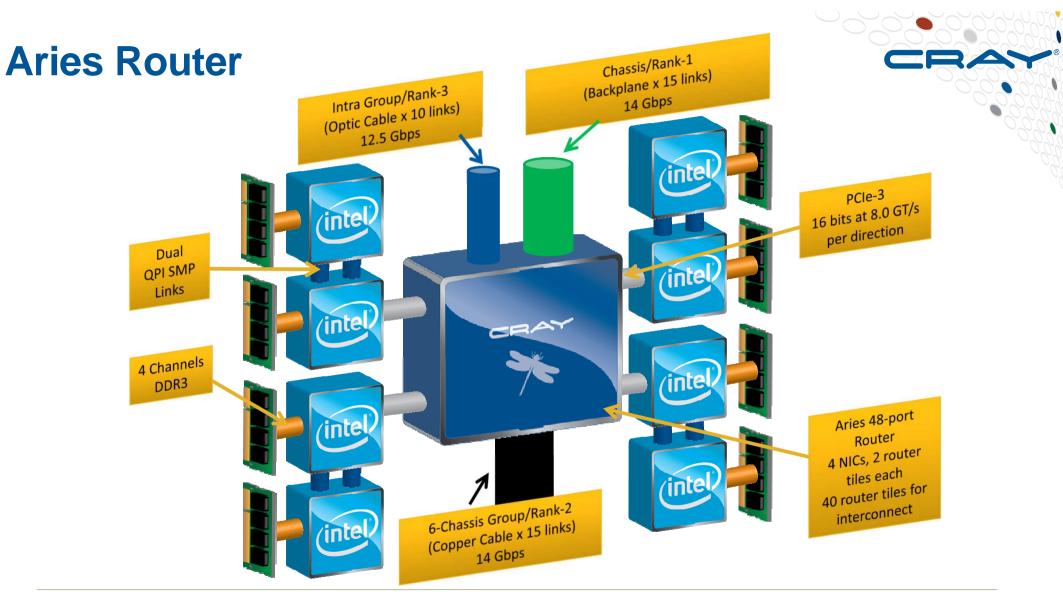
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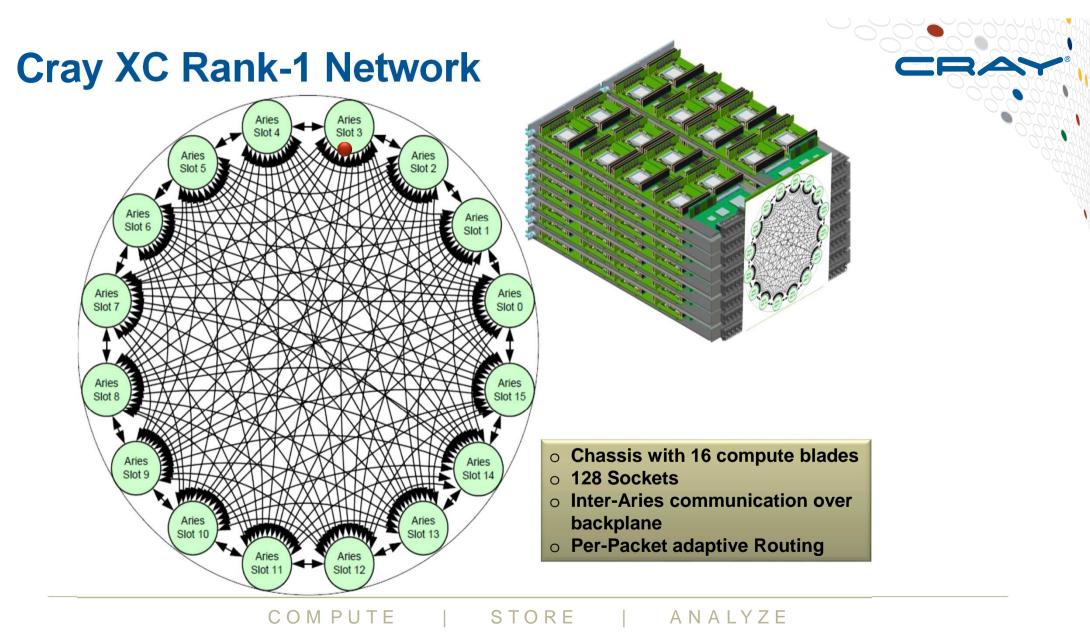
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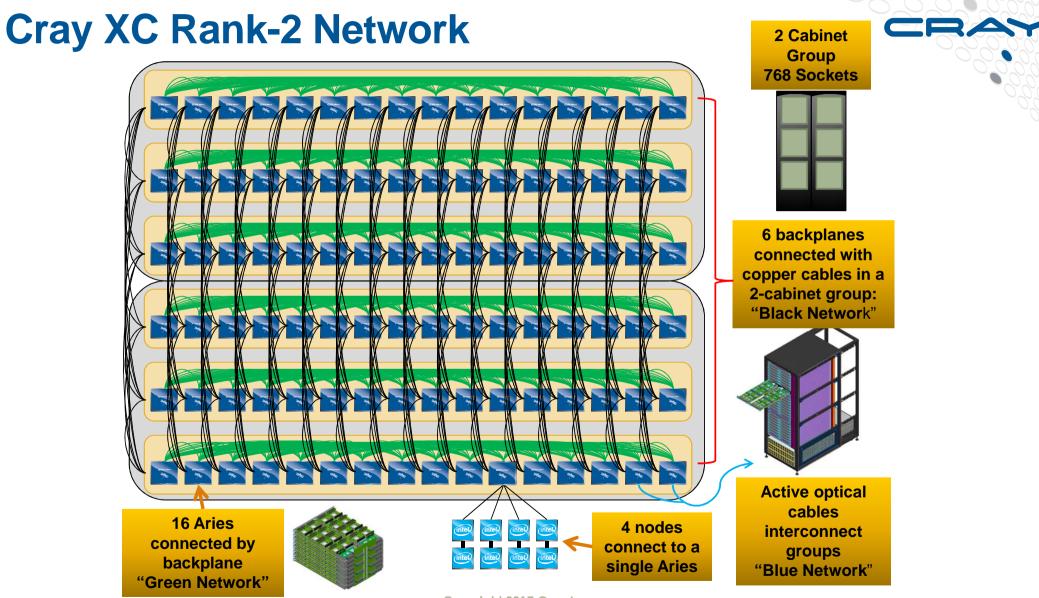
Aries Network Infrastructure



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Cray XC40 Rank-2 Network

- Cray XC40 two-cabinet group
 - 768 Sockets
 - 96 Aries Chips
- All copper and backplanes signals running at 14 Gbps





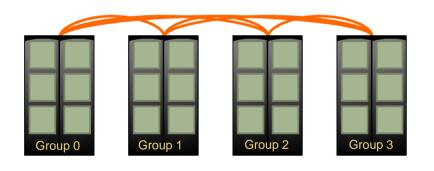
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Cray XC40 Rank-3 Network

- An all-to-all pattern is wired between the groups using optical cables (blue network)
- Up to 240 ports are available per 2-cabinet group
- The global bandwidth can be tuned by varying the number of optical cables in the group-to-group connections



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Example: A 4-group system is interconnected with 6 optical "bundles". The "bundles" can be configured between 20 and 80 cables wide

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Scalability Metrics for Supercomputing

Standard metrics

- Infrastructure
- Compute
- Power
- Performance of collective operations across system
- MPI stack memory footprint
- Job start-up times

Why Dragonfly?

• Cost

- Dragonfly minimizes the use of active optical components
- Eliminates need for Director switches

Scalability

- Topology scales to very large systems
- Flat average hop count and latency

Simplicity

- Implemented without external switches
- No HBAs or separate NICs and Routers

Performance

• More than just a case of clever wiring, this topology leverages state-ofthe-art adaptive routing that Cray developed with Stanford University

Why Dragonfly?

Comparison with Fat-Tree

- Cost
 - Fat-Tree increases cost per node with system size
 - 2x optical links for the same global bandwidth
 - Requires external ToR and Director class routers
- Latency
 - Fat-Tree requires 2 optical hops per route vs. 1 for Dragonfly
- Load balancing
 - Application traffic patterns (all-to-all, uniform random) self load-balance using Dragonfly
 - Large electrical group captures most of local load

Infrastructure Topology

Dragonfly average hop-count is low and flat

- Up to 2 hops within source group
- One hop to destination group
- Up to 2 hops within destination group
- Hop-count stays low out to very big networks

Per-node bisectional bandwidth is flat

- Half of all-to-all connected global links
- Grows linearly with system size out to very big networks

Power Consumption – Energy to Solution

Energy to solution = time to solution x system power x power utilization efficiency

Time to solution

- Reducing time is the most effective element
- Improvements in scalability directly affect time to solution

System power

- XC system rack consumes ~80kW under load
- 420W per dual socket node
- 15W per node Aries Dragonfly
- 48W per node Fat-Tree
- Reduces consumption by 625kWH \$3.5m running costs

• Power utilization efficiency (PUE)

- 480V distribution, 48V rack, ~1V component
- Warm water cooling
- Cray XC is typically 1.1 1.25 PUE
- Google data center 1.21 PUE
- Microsoft data center 1.22 PUE

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PUE = Facility Energy / System Energy

Collective Operations

• Aries NICs includes a Collective Offload Engine

- No CPU involvement
- Latency optimized
- Scales with network size
- Up to radix 32
- 32 / 64 bit integer and floating point add
- Min/max, compare & swap, bit operations
- No topological tree dependency

• High branching ratio = shallow trees

- Low latency
- Require only 3 stage reduction tree for common operations on Trinity

Collective Operations Offload

Typical 2 phase operation

- Ready phase
 - Leaf Nodes join the reduction
 - Reduction operator applied as data moves towards the Root Node

Parent

.....

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Leaf

Leaf

Leaf

Parent

.....

Leaf

Root

Leaf

Leaf

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Leaf

Parent

Leaf

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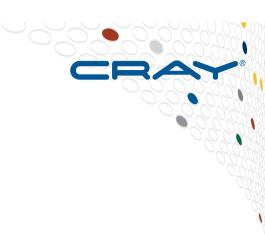
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- Multicast phase
 - Results pass from Root to Leaf nodes

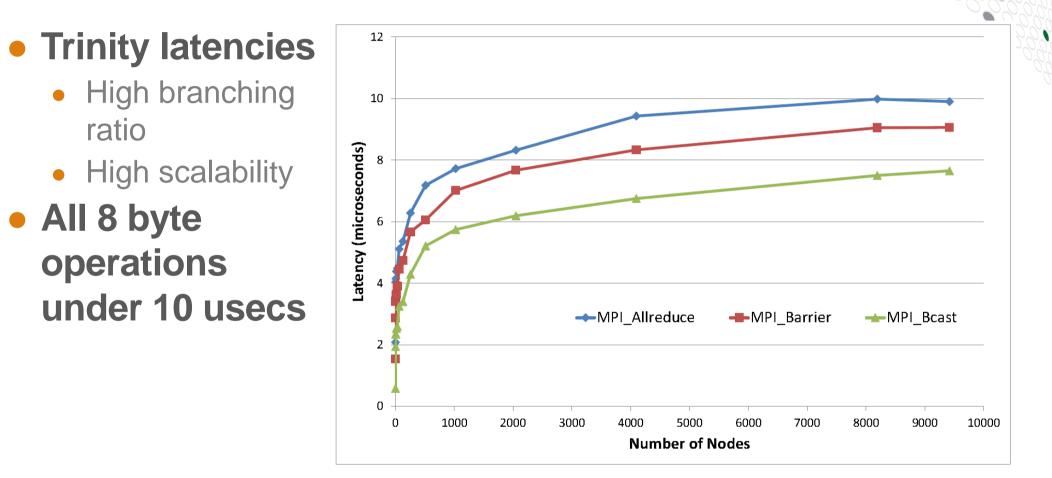
Operations are non-blocking

• 128 per job for a radix-32 tree

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Collective Operations Offload



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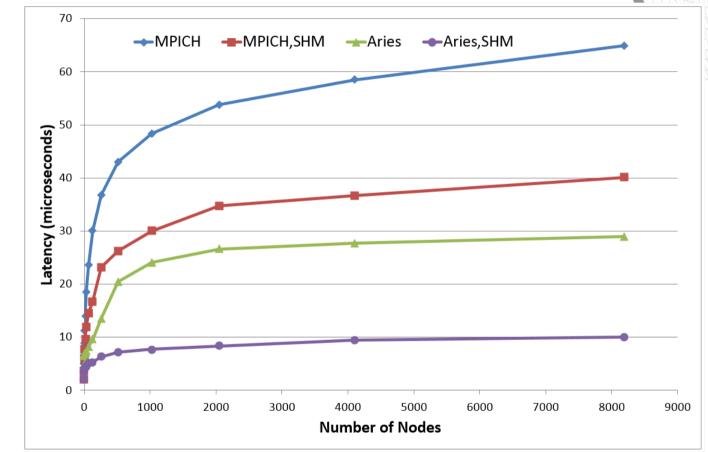
Collective Operations – Local shared-memory

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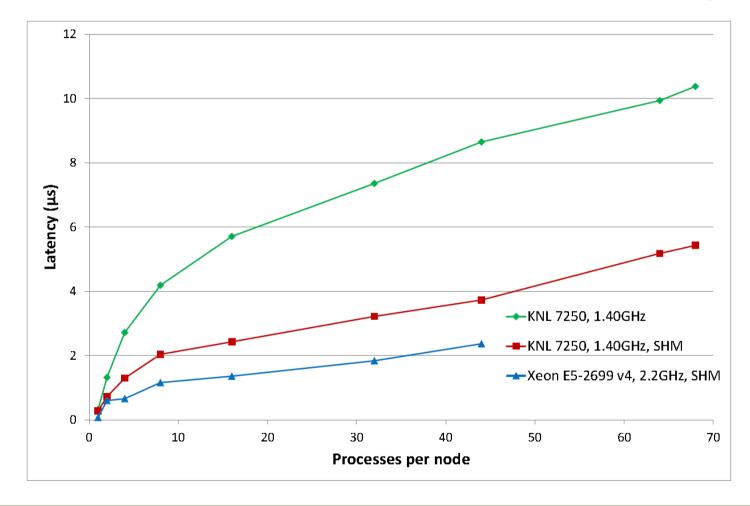
- MPI_Allreduce
- Offload changes the balance
 - Now intra-node reduction becomes significant
 - Especially on many-core nodes

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Collective Operations – Local shared-memory



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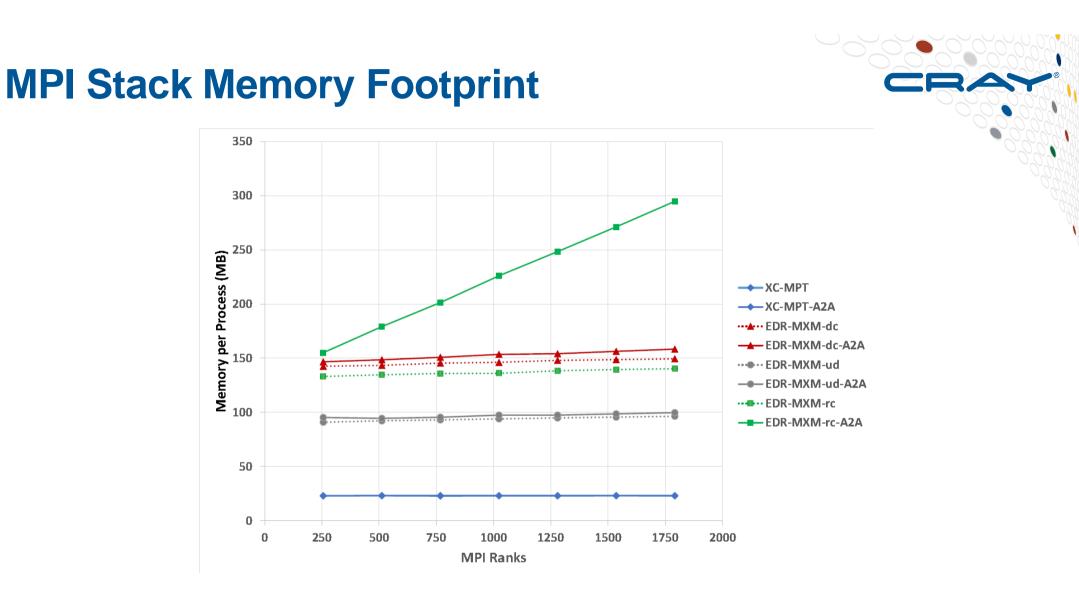
MPI Stack Memory Footprint

Memory used by MPI stack presents scalability barrier

- Each process maintains static state for every other process
- Total memory used for inter-processor comms includes:
 - MPI Virtual Channel structures peer-to-peer state
 - Process Management Interface
 - Transport Layer memory
 - Per-node shared state

• Cray MPI implements dynamic MPI Virtual Channels

- Significantly reduces MPI stack memory footprint
- Memory allocated only when ranks contact each other
- Utilizes connectionless RMAs no VC usage



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Job Startup Times

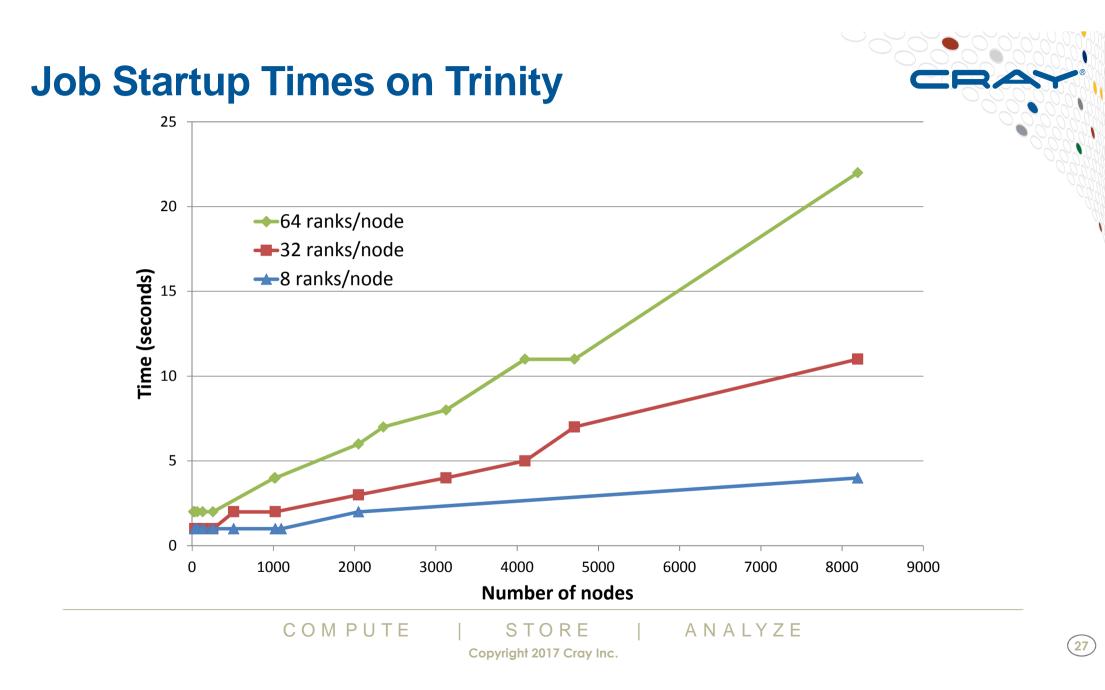
• Startup times rise with job size

- Reduced efficiency for large jobs
- Decreased system utilization

• Cray XC startup times illustrated

- 301,248 process job, 32 ranks/node 12 secs
- 64 ranks/node 24 secs

Dynamic libraries must provision I/O resources to ensure fast loading



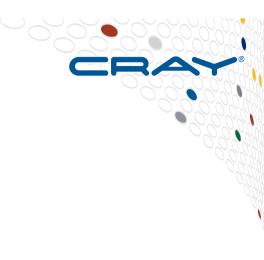
Conclusions



- Key aspects of scalability identified
 - Infrastructure, MPI stack, collectives & cost
- Investigated effects at scale using Trinity
 - Running >300,000 MPI ranks

Excellent performance demonstrated in key metrics

- Benchmarks, acceptance tests & early application use
- Interconnect & its software are critical factors
 - Dragonfly network & software stack are key elements of Trinity
 - Results shown could not have been achieved otherwise



Thank You

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